



Joinads

A retargetable control-flow construct for reactive, parallel and concurrent programming

Tomáš Petříček (tomas.petricek@cl.cam.ac.uk)

University of Cambridge, UK

Don Syme (don.syme@microsoft.com)

Microsoft Research, Cambridge, UK

The two key points of the talk

■ Language extension

We add language support for concurrent, parallel and reactive programming

■ Multi purpose

We do this without committing the language to one particular programming model

■ We extend F# computation expressions

Similar approach could be used in other functional languages (especially Haskell's do-notation)

Reactive, concurrent and parallel

□ Programming with *futures*

- Running in background and eventually gives a result
- Language support in Manticore (Fluet et al. 2008)

□ Event-based programming

- Lightweight threads, communicating using *events*
- Functional Reactive Programming (Elliott 2000)

□ Join-calculus

- *Joins* execute when certain *channels* contain values
- Both languages (Conchon, Fessant 1999) and libraries (Russo 2007)

Bringing programming models to practice

- Language-based solutions

- Language supports only one model

- Library-based encodings

- Restricted syntax is limiting



- **Our approach:** Support a recurring pattern

- Successfully used by monads (and arrows & idioms)
 - One syntactic extension works for many libraries

Overview

▣ Background

Computation expressions overview

▣ Our extension

Choosing between computations

Merging computations

What are joinads?

▣ Interesting relations

Joinads and other computation types

Computation expressions by example

Function of type
int -> Event<int>

Computation expression
creating **Event<int>**

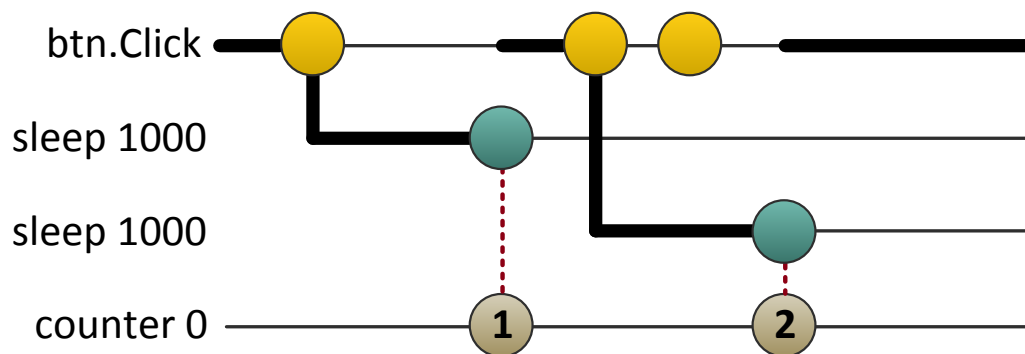
Recursive
looping

```
let rec counter n = event {  
  let! args = btn.Click  
  let! time = Event.sleep 1000  
  return n + 1  
  return! counter (n + 1) }
```

Waiting for the first
occurrence of an event
Event<MouseInfo>

Trigger the created event

■ Event is modeled as a sequence of time-value pairs



F# computation expressions

□ Computation expression syntax

cexpr = **let** *pat* = *expr* **in** *cexpr* Binding value
 | **let!** *pat* = *expr* **in** *cexpr* Binding computation
 | **return** *expr* Returning value
 | **return!** *expr* Returning computation
 | **match** *expr-list* **with** ... Pattern matching on values

□ Notation for writing computations ('do' in Haskell)

□ Translates to primitive function calls

bind : $M\langle a \rangle \rightarrow (a \rightarrow M\langle b \rangle) \rightarrow M\langle b \rangle$
unit : $a \rightarrow M\langle a \rangle$
combine : $M\langle a \rangle \rightarrow M\langle a \rangle \rightarrow M\langle a \rangle$

F# computation expressions

□ Computation expression syntax

<i>cexpr</i> =	let <i>pat</i> = <i>expr</i> in <i>cexpr</i>	Binding value
	let! <i>pat</i> = <i>expr</i> in <i>cexpr</i>	Binding computation
	return <i>expr</i>	Returning value
	return! <i>expr</i>	Returning computation
	match <i>expr-list</i> with ...	Pattern matching on values

□ Our extension adds the obvious

| **match!** *expr-list* **with** ... Pattern matching on computations

□ ...and two primitive functions for the translation

□ They specify what **match!** actually means

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Joinads and other computation types

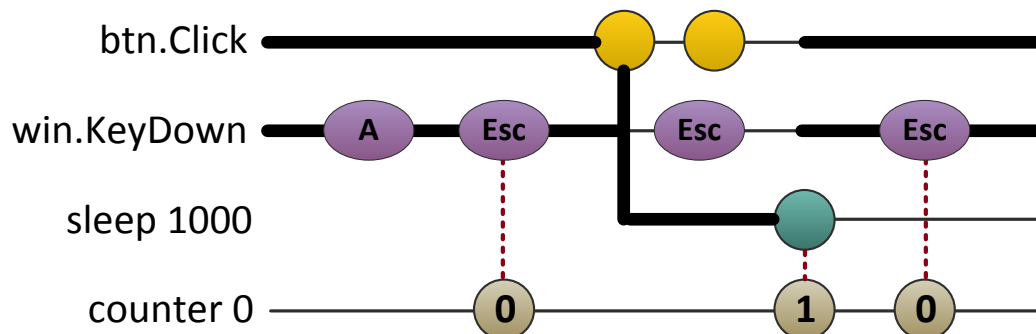
Choosing between computations

- Operation **choose** composes multiple clauses
 - Wait for events in parallel & run the first enabled body

Button was clicked
Wait & increment

Runs only when key
matches pattern

```
let rec counter n = event {  
  match! btn.Click, win.KeyDown with  
  | !_, _ -> let! _ = Event.sleep 1000  
              return n + 1  
              return! counter (n + 1)  
  | _, !Esc -> return 0  
              return! counter 0 }
```



What patterns can we write?

■ New syntactic category *computation pattern*

$cexpr = \mathbf{match!} \text{ } expr\text{-list} \mathbf{with}$ Pattern matching on computations
 $cpat\text{-list} \rightarrow cexpr \mid \dots$ with a list of clauses

$cpat = _$ Ignore computation pattern
 $\mid !pat$ Bind computation using standard pattern

■ Note the difference between “ $_$ ” and “ $! _$ ”

- **!Esc** is a *non-exhaustive* computation binding
- **!_** is *exhaustive* but needs a value to match on
- **_** matches even if we don't have a value

Merging computations

■ Binding values from multiple computations

- All clauses so far had only single binding pattern
- Operation **merge** combines computations

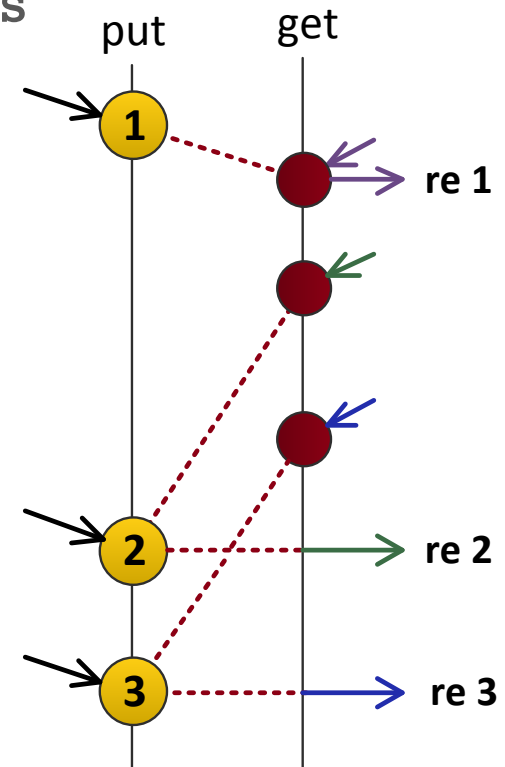
Asynchronous input &
Synchronous output

```
let put = new Channel<int>()  
let get = new Channel<ReplyChan<string>>()
```

```
let buffer = join {  
  match! put, get with  
  | !num, !chnl ->  
    reply chnl (sprintf "re %d" num)
```

Computation
based on joins

Pattern “joining”
the two channels



What is a joinad?

```
map      : (a → b) → M<a> → M<b>
merge    : M<a> → M<b> → M<a * b>
choose   : list<M<option<M<a>>>> → M<a>
```

- The **match!** syntax translates to these
 - **merge** – Combines two computations into a single
 - **choose** – Finds the first enabled computation from a list of clauses and returns computation that runs the body
- **Call to Action: Formalization of Joinads**
 - Are these the simplest primitives we can use?
 - How to find complete laws about the primitives?

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Joinads and other computation types

Joinads and monads

- Joinads do not imply monads or otherwise
 - Many computations are both joinad and monad
- Can we get **merge** inside monad for free?
 - The type is $M\langle a \rangle \rightarrow M\langle b \rangle \rightarrow M\langle a * b \rangle$
 - Want commutativity $\text{merge } u \ v \equiv \text{map swap } (\text{merge } v \ u)$

```
let merge ma mb = m {  
  let! a = ma  
  let! b = mb  
  return a, b }
```

```
let merge ma mb = m {  
  let! b = mb  
  let! a = ma  
  return a, b }
```

**Commutative
monads!**

Summary & Questions?

■ Language extension for multiple models

- **Reactive** based on events (similar to FRP)
- **Parallel** based on futures (related to Manticore)
- **Concurrent** based on join calculus (JoCaml, C_ω)
- ...and possibly many others

■ Theoretically interesting

- More work to be done on the formal model...

tomas.petricek@cl.cam.ac.uk



The end of the universe

Joined computations for futures

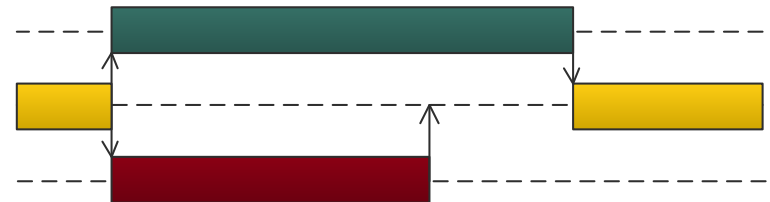
- Future is computation running in background
 - Binding means waiting for the completion

```
let multiply f1 f2 = future {  
  match! f1, f2 with  
  | !a, !b -> return a * b  
  | !0, _   -> return 0  
  | _, !0   -> return 0 }
```

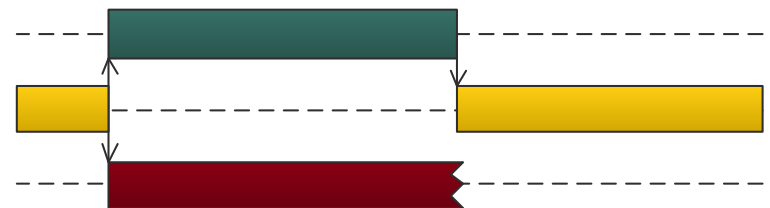
One of the futures
completed and
produced zero

Both futures
completed

Case !a, !b



Case !0, _



Desugaring of computation expressions

```
let rec counter n = event {  
  let! _ = btn.Click  
  let! _ = Event.sleep 1000  
  return n + 1  
  return! counter (n + 1) }
```

Waiting for an event

■ Functions are associated with the **event** builder

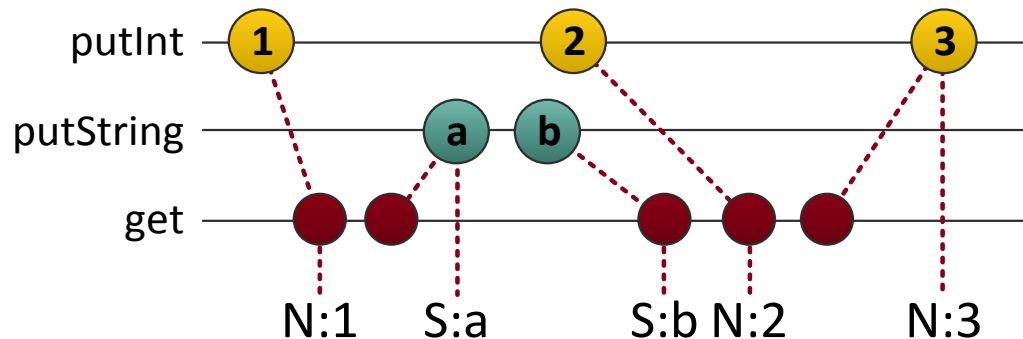
■ **return** and **let!** translate to **Return** and **Bind**

■ Sequencing of expressions translates to **Combine**

```
let rec counter n =  
  event.Bind(btn.Click, fun _ ->  
    event.Bind(Event.sleep 1000, fun _ ->  
      event.Combine  
        ( event.Return(n + 1),  
          counter (n + 1) )))
```

Continuation called
once when event occurs

Desugaring of joinads



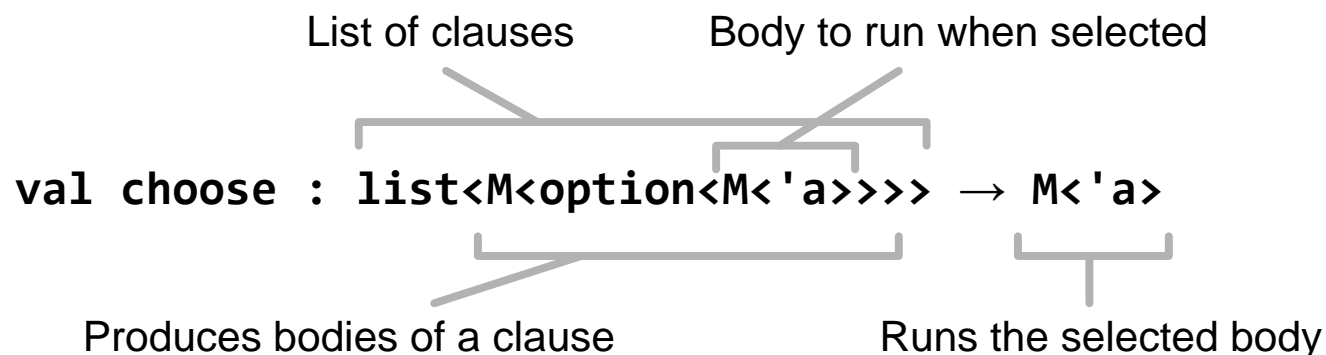
```
let putInt = new Channel<int>()
let putString = new Channel<string>()
let get = new Channel<ReplyChannel<string>>()
```

```
let buffer =
  join.Choose
    [ join.Merge(get, putInt) |> join.Map (fun (chnl, n) ->
      join { chnl.Reply("Number: " + n.ToString()) });
      join.Merge(get, putString) |> join.Map (fun (chnl, s) ->
        join { chnl.Reply("String:" + s) }) ]
```

Pattern matching
done inside **Map**

Multiple binding patterns
turned into **Merge**

Choose operation explained



■ Type signature resembles monadic **join**

- Should behave the same for singleton list with “Some”
- Outer computation
 - Maps matching inputs into clauses to be executed
- Inner computation
 - Represents the body

Joinad laws: Where do they come from?

□ Transformations that shouldn't change meaning

match! m **with** $!var \rightarrow expr$ \equiv **let!** $var = m$ **in** $expr$

Replace trivial
“match!” with binding

match! m { **return** e_1 },
 m { **return** e_2 } **with** \equiv **match** e_1, e_2 **with**
 $| !var_1, !var_2 \rightarrow cexpr$ $| var_1, var_2 \rightarrow cexpr$

Pattern matching
on two “units”

match! $\dots, m_{p(i)}, \dots$ **with**
 $| \dots, cpat_{1,p(i)}, \dots \rightarrow cexpr_1$ \equiv **match!** \dots, m_i, \dots **with**
 $| \dots, cpat_{1,i}, \dots \rightarrow cexpr_1$
 $| \dots$ $| \dots$
 $| \dots, cpat_{k,p(i)}, \dots \rightarrow cexpr_k$ $| \dots, cpat_{k,i}, \dots \rightarrow cexpr_k$

Reordering of
computations
& patterns

match! m **with**
 $| !var_1 \rightarrow \langle cexpr \rangle_1$ \equiv **match!** m **with**
 $| !var_1 \rightarrow \langle cexpr \rangle_1$
 $| !var_2 \rightarrow \langle cexpr \rangle_2$

Match first enabled clause

Joinad laws: Simplified form

■ Merge operation (written as \mathbb{II})

■ Commutativity is related to commutative monads

$$u \mathbb{II} (v \mathbb{II} w) \equiv \text{map assoc } ((u \mathbb{II} v) \mathbb{II} w) \quad (\textit{associativity})$$

$$u \mathbb{II} v \equiv \text{map swap } (v \mathbb{II} u) \quad (\textit{commutativity})$$

$$\text{unit } (u, v) \equiv (\text{unit } u) \mathbb{II} (\text{unit } v) \quad (\textit{unit merge})$$

where $\text{assoc } ((a, b), c) = (a, (b, c))$ **and** $\text{swap } (a, b) = (b, a)$

■ Choose operation

■ Should always select the first enabled clause (formal definition doesn't make things much simpler)

■ For monads, should generalize **bind** operation

Translation of Joinads

- **Merge** inputs for pattern matching and **map**
- Translate clauses using $\langle - \rangle$ and apply **choose**

$$\begin{aligned} \ll \text{match! } \text{expr}_1, \dots, \text{expr}_k \text{ with } ccl_1 \mid \dots \mid ccl_p \gg_m &\equiv \\ \text{let } v_1 = \text{expr}_1 \text{ in } \dots \text{ let } v_k = \text{expr}_k \text{ in} & \\ \text{choose}_m [\langle ccl_1 \rangle_m, (v_1, \dots, v_k); \dots ; \langle ccl_p \rangle_m, (v_1, \dots, v_k)] & \end{aligned}$$

$$\begin{aligned} \langle cpat_1, \dots, cpat_k \rightarrow cexpr \rangle_{m, (v_1, \dots, v_k)} &\equiv \\ \text{map}_m (\text{function } (pat_1, \dots), pat_n \rightarrow \text{Some } \ll cexpr \gg_m & \\ \mid _ \rightarrow \text{None}) \text{cargs} & \end{aligned}$$

$$\begin{aligned} \text{where } \{ (pat_1, v_1), \dots, (pat_n, v_n) \} &= \{ (pat, v_i) \mid cpat_i = !pat, 1 \leq i \leq k \} \\ \text{cargs} &= v_1 \oplus_m \dots \oplus_m v_{n-1} \oplus_m v_n \quad \text{for } n \geq 1 \end{aligned}$$